

## Matrix Sieve – New Algorithm for Finding Prime Numbers

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### Abstract

A new deterministic sieving algorithm based on derived “matrix definition” of prime numbers is proposed. The algorithm allows to calculate indexes  $P$  of prime numbers in two sequences:  $S_1(P) = 5 + 6P = 5, 11, 17, \dots; P = 0, 1, 2, 3, \dots$  and  $S_2(P) = 7 + 6P = 7, 13, 19, \dots; P = 0, 1, 2, 3, \dots$  in a given range of natural numbers  $(N_1, N_2)$ . Also general primality criteria and twin-prime criteria are formulated. C++ program for finding primes in given range  $(N_1, N_2)$  and C++ program for primality testing of given natural number  $N$  are presented in Attachments 1 and 2.

### 1. Matrix general expression for composite numbers

Consider a sequence of natural numbers from which members divisible by 2 and 3 are removed

$$S(p) = 5, 7, 11, 13, 17, 19, 23, 25, 29, \dots = \begin{cases} 3p + 5, p = 0, 2, 4, 6, 8, \dots \\ 3(p - 1) + 7, p = 1, 3, 5, 7, \dots \end{cases} \quad (1)$$

Sequence  $S(p)$  can be divided into two sequences  $S_1(i)$  and  $S_2(i)$ :

$$S_1(i) = 5 + 6i = 5, 11, 17, \dots; i = 0, 1, 2, 3, \dots \quad (2)$$

$$S_2(i) = 7 + 6i = 7, 13, 19, \dots; i = 0, 1, 2, 3, \dots \quad (3)$$

Sequence  $S(p)$  contains all primes (except 2 and 3) and, using definitions of  $S_1(i)$  and  $S_2(i)$ , all composite numbers (except divisible by 2 and 3) of four different types:

$$FF(i, j) = S_1(i) * S_1(j) = (5 + 6i)(5 + 6j); \quad SS(i, j) = S_2(i) * S_2(j) \\ = (7 + 6i)(7 + 6j);$$

$$SF(i, j) = S_1(i) * S_2(j) = (5 + 6i)(7 + 6j); \quad FS(i, j) = S_1(j) * S_2(i) = (5 + 6j)(7 + 6i);$$

which constitute four 2-dimensional arrays:

$$FF(i, j) = 25 + 30(i + j) + 36ij = \begin{pmatrix} 25 & 55 & 85 & 115 & 145 & \dots \\ 55 & 121 & 187 & 253 & 319 & \dots \\ 85 & 187 & 289 & 391 & 493 & \dots \\ 115 & 253 & 391 & 529 & 667 & \dots \\ 145 & 319 & 493 & 667 & 841 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix}; i, j =$$

0,1,2,3,4 ....(4)

$$SS(i, j) = 49 + 42(i + j) + 36ij = \begin{pmatrix} 49 & 91 & 133 & 175 & 217 & \dots \\ 91 & 169 & 247 & 325 & 403 & \dots \\ 133 & 247 & 361 & 475 & 589 & \dots \\ 175 & 325 & 475 & 625 & 775 & \dots \\ 217 & 403 & 589 & 775 & 961 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix}; i, j =$$

0,1,2,3,4 .... (5)

$$FS(i, j) = 35 + 6(5i + 7j) + 36ij = \begin{pmatrix} 35 & 77 & 119 & 161 & 203 & \dots \\ 65 & 143 & 221 & 299 & 377 & \dots \\ 95 & 209 & 323 & 437 & 551 & \dots \\ 125 & 275 & 425 & 575 & 725 & \dots \\ 155 & 341 & 527 & 713 & 899 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix}; i, j =$$

0,1,2,3,4 ....(6)

$$SF(i, j) = 35 + 6(7i + 5j) + 36ij = \begin{pmatrix} 35 & 65 & 95 & 125 & 155 & \dots \\ 77 & 143 & 209 & 275 & 341 & \dots \\ 119 & 221 & 323 & 425 & 527 & \dots \\ 161 & 299 & 437 & 575 & 713 & \dots \\ 203 & 377 & 551 & 725 & 899 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix}; i, j =$$

0,1,2,3,4 .... (7)

Substituting numbers in the above matrices with their corresponding indexes  $p$  in accordance with equation (1), the following four matrices can be obtained:

$$ff(i, j) = \begin{pmatrix} 7 & 17 & 27 & 37 & \dots & \dots \\ 17 & 39 & 61 & 83 & \dots & \dots \\ 27 & 61 & 95 & 129 & \dots & \dots \\ 37 & 83 & 129 & 175 & \dots & \dots \\ 47 & 105 & 163 & 221 & \dots & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \end{pmatrix} \quad (8)$$

$$ss(i, j) = \begin{pmatrix} 15 & 29 & 43 & 57 & \dots & \dots \\ 29 & 55 & 81 & 107 & \dots & \dots \\ 43 & 81 & 119 & 157 & \dots & \dots \\ 57 & 107 & 157 & 207 & \dots & \dots \\ 71 & 133 & 195 & 257 & \dots & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \end{pmatrix} \quad (9)$$

$$fs(i, j) = \begin{pmatrix} 10 & 24 & 38 & 52 & \dots & \dots \\ 20 & 46 & 72 & 98 & \dots & \dots \\ 30 & 68 & 106 & 144 & \dots & \dots \\ 40 & 90 & 140 & 190 & \dots & \dots \\ 50 & 112 & 174 & 236 & \dots & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \end{pmatrix} \quad (10)$$

$$sf(i, j) = \begin{pmatrix} 10 & 20 & 30 & 40 & \dots & \dots \\ 24 & 46 & 68 & 90 & \dots & \dots \\ 38 & 72 & 106 & 140 & \dots & \dots \\ 52 & 98 & 144 & 190 & \dots & \dots \\ 66 & 124 & 182 & 240 & \dots & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \end{pmatrix} \quad (11)$$

In general form these arrays can be expressed as

$$ff(i, j) = 12ij - 2(i + j) - 1; i, j = 1, 2, 3, 4 \dots \quad (12)$$

$$ss(i, j) = 12ij + 2(i + j) - 1; i, j = 1, 2, 3, 4 \dots \quad (13)$$

$$fs(i, j) = 12ij - 2(i - j) - 2; i, j = 1, 2, 3, 4 \dots \quad (14)$$

$$sf(i, j) = 12ij + 2(i - j) - 2; i, j = 1, 2, 3, 4 \dots \quad (15)$$

## 2. “Matrix definition” of prime numbers

As it follows from expressions (12)-(15), arrays  $ff(i, j)$  and  $ss(i, j)$  are symmetric and are comprised of odd integers (i.e. odd indexes of elements of  $S(p)$ ); arrays  $sf(i, j)$  and  $fs(i, j)$  are transposes of each other and are comprised of even integers (i.e. even indexes of elements of  $S(p)$ ). Thus, all elements of arrays  $ff(i, j)$  and  $ss(i, j)$  are indexes of members of  $S_2(p)$  and all elements of arrays  $sf(i, j)$  and  $fs(i, j)$  are indexes of members of  $S_1(p)$ . It means that all composite members of the sequence  $S_1(p)$  correspond to formula

$$S_1(P) = 6P + 5 = (5 + 6i)(7 + 6j); i, j = 0, 1, 2, 3, 4 \dots$$

And all composite members of the sequence  $S_2(p)$  correspond to formulae

$$S_2(P) = 6P + 7 = (5 + 6i)(5 + 6j); i = 0, 1, 2, 3, 4 \dots, j \geq i \text{ or}$$

$$S_2(P) = 6P + 7 = (7 + 6i)(7 + 6j); i = 0, 1, 2, 3, 4 \dots, j \geq i$$

. Let us now consider sequences  $S_1(P)$  and  $S_2(P)$  separately as defined below

$$S_1(P) = 6P + 5; P = 0, 1, 2, 3, \dots = p/2 \quad (16)$$

$$S_2(P) = 6P + 7; P = 0, 1, 2, 3, \dots = (p - 1)/2 \quad (17)$$

Indexes  $P_1(i, j)$ ,  $P_2(i, j)$  and  $P_3(i, j)$ ,  $P_4(i, j)$  corresponding to composite numbers in  $S_1(P)$  and  $S_2(P)$  respectively are then defined as:

for  $S_1(P)$

$$P_1(i, j) = 6ij - i + j - 1 = \begin{pmatrix} 5 & 12 & 19 & 26 & 33 & 40 & \dots \\ 0 & 23 & 36 & 49 & 62 & 75 & \dots \\ 0 & 0 & 53 & 72 & 91 & 110 & \dots \\ 0 & 0 & 0 & 95 & 120 & 145 & \dots \\ 0 & 0 & 0 & 0 & 149 & 180 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 215 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i \quad (18)$$

$$P_2(i, j) = 6ij + i - j - 1 = \begin{pmatrix} 0 & 10 & 15 & 20 & 25 & 30 & \dots \\ 0 & 0 & 34 & 45 & 56 & 67 & \dots \\ 0 & 0 & 0 & 70 & 87 & 104 & \dots \\ 0 & 0 & 0 & 0 & 118 & 141 & \dots \\ 0 & 0 & 0 & 0 & 0 & 178 & \dots \\ \cdot & \cdot & \cdot & \cdot & \cdot & 0 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i + 1 \quad (19)$$

for  $S_2(P)$

$$P_3(i, j) = 6ij - i - j - 1 = \begin{pmatrix} 3 & 8 & 13 & 18 & 23 & 28 & \dots \\ 0 & 19 & 30 & 41 & 52 & 63 & \dots \\ 0 & 0 & 47 & 64 & 81 & 98 & \dots \\ 0 & 0 & 0 & 87 & 110 & 133 & \dots \\ 0 & 0 & 0 & 0 & 139 & 168 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 203 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i \quad (20)$$

$$P_4(i, j) = 6ij + i + j - 1 = \begin{pmatrix} 7 & 14 & 21 & 28 & 35 & 42 & \dots \\ 0 & 27 & 40 & 53 & 66 & 79 & \dots \\ 0 & 0 & 59 & 78 & 97 & 116 & \dots \\ 0 & 0 & 0 & 103 & 128 & 153 & \dots \\ 0 & 0 & 0 & 0 & 159 & 190 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 227 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i \quad (21)$$

Thus, we have obtained a matrix representation of indexes of all composite numbers, which allows us to determine indexes of all prime numbers; so “**matrix definition**” of **prime numbers**” can be formulated as follows:

Natural numbers not contained in arrays  $P_1(i, j)$  and  $P_2(i, j)$  are indexes  $P$  of all prime numbers in sequence  $S_1(P)$  and natural numbers not contained in arrays  $P_3(i, j)$  and  $P_4(i, j)$  are indexes  $P$  of all prime numbers in sequence  $S_2(P)$ .

Additionally, equations (18) – (21) provide algebraic relationships connecting integers  $i$  and  $j$  with indexes of composite numbers in  $S_1(P)$  and  $S_2(P)$ . Now, for sieving on an interval purposes, we need to establish a relationship between our range of interest  $(N_1, N_2)$  and the corresponding area of change of integers  $i$  and  $j$ :

Since  $j \geq i$ , maximum value of  $i$  will be when  $i = j$  and it approximately equals

$$i_{max} = \frac{\sqrt{N_2}}{6} \quad (22)$$

$$\text{For array } P_1(i, j) \text{ we have: } j_{min} = \frac{\frac{N_1+i+1}{6}}{6i+1}; \quad j_{max} = \frac{\frac{N_2+i+1}{6}}{6i+1}; \quad (23)$$

$$\text{For array } P_2(i, j): \quad j_{min} = \frac{\frac{N_1-i+1}{6}}{6i-1}; \quad j_{max} = \frac{\frac{N_2-i+1}{6}}{6i-1}; \quad (24)$$

$$\text{For array } P_3(i, j): \quad j_{min} = \frac{\frac{N_1+i+1}{6}}{6i-1}; \quad j_{max} = \frac{\frac{N_2+i+1}{6}}{6i-1}; \quad (25)$$

$$\text{For array } P_4(i, j): \quad j_{min} = \frac{\frac{N_1-i+1}{6}}{6i+1}; \quad j_{max} = \frac{\frac{N_2-i+1}{6}}{6i+1}; \quad (26)$$

**Matrix sieving algorithm** can be formulated as follows:

In order to find all prime numbers in the range from  $N_1$  to  $N_2$  it is necessarily to remove from sequence  $S_1(P)$  members with indexes  $P_1(i, j)$ ,  $P_2(i, j)$ :

$$S_1(P) = 6P + 5; \quad P = 0, 1, 2, 3, \dots$$

$$P_1(i, j) = 6ij - i + j - 1; \quad i = 1, 2, \dots, \left\lfloor \frac{\sqrt{N_2}}{6} \right\rfloor; \quad j = \left\lfloor \frac{\frac{N_1+i+1}{6}}{6i+1} \right\rfloor, \left\lfloor \frac{\frac{N_1+i+1}{6}}{6i+1} \right\rfloor + 1, \dots, \left\lfloor \frac{\frac{N_2+i+1}{6}}{6i+1} \right\rfloor; \quad j \geq i \quad (27)$$

$$P_2(i, j) = 6ij + i - j - 1; \quad i = 1, 2, \dots, \left\lfloor \frac{\sqrt{N_2}}{6} \right\rfloor; \quad j = \left\lfloor \frac{\frac{N_1-i+1}{6}}{6i-1} \right\rfloor, \left\lfloor \frac{\frac{N_1-i+1}{6}}{6i-1} \right\rfloor + 1, \dots, \left\lfloor \frac{\frac{N_2-i+1}{6}}{6i-1} \right\rfloor; \quad j \geq i+1 \quad (28)$$

And remove from sequence  $S_2(P)$  members with indexes  $P_3(i, j)$ ,  $P_4(i, j)$ :

$$S_2(P) = 6P + 7; \quad P = 0, 1, 2, 3, \dots$$

$$P_3(i, j) = 6ij - i - j - 1; \quad i = 1, 2, \dots, \left\lfloor \frac{\sqrt{N_2}}{6} \right\rfloor; \quad j = \left\lfloor \frac{\frac{N_1+i+1}{6}}{6i-1} \right\rfloor, \left\lfloor \frac{\frac{N_1+i+1}{6}}{6i-1} \right\rfloor + 1, \dots, \left\lfloor \frac{\frac{N_2+i+1}{6}}{6i-1} \right\rfloor; \quad j \geq i \quad (29)$$

$$P_4(i, j) = 6ij + i + j - 1; \quad i = 1, 2, \dots, \left\lfloor \frac{\sqrt{N_2}}{6} \right\rfloor; \quad j = \left\lfloor \frac{\frac{N_1-i+1}{6}}{6i+1} \right\rfloor, \left\lfloor \frac{\frac{N_1-i+1}{6}}{6i+1} \right\rfloor + 1, \dots, \left\lfloor \frac{\frac{N_2-i+1}{6}}{6i+1} \right\rfloor; \quad j \geq i \quad (30)$$

### 3. C++ program based on matrix sieve algorithm

In programming code were used following notations:

$$pr1 = P_{\min}; pr2 = P_{\max};$$

R1[q], R2[r], S1[q], S2[r]-additional arrays corresponding to the range of P (pr1;pr2).

q- index of the arrays R1[q], S1[q].

r – index of the arrays R2[r], S2[r].

$$i2 = imax; j1 = jmin; j2 = jmax;$$

$$P1 = P_1(i, j); P2 = P_2(i, j); P3 = P_3(i, j); P4 = P_4(i, j);$$

Expressions (18)-(21) for programming code can be rewritten as:

$$P1[i;j] = 5 + 5*(i-1) + (7 + 6*(i-1))*(j-1) \quad = (18a)$$

$$P2[i;j] = 5 + 7*(i-1) + (5 + 6*(i-1))*(j-1) \quad (19a)$$

$$P3[i;j] = 3 + 5*(i-1) + (5 + 6*(i-1))*(j-1) \quad (20a)$$

$$P4[i;j] = 7 + 7*(i-1) + (7 + 6*(i-1))*(j-1) \quad (21a)$$

Using above equations (18a)-(21a) with area of change of i and j determined as before (22)-(26), the C++ program was developed to find prime numbers in the range ( $N_1; N_2$ ) (Appendix 1). Program was successfully tested on the ordinary notebook up to  $N = 2\,000\,000\,000\,000\,000\,000\,000$ , for this value run time equals 10 seconds.

### 4. Primality criteria

Taking into account above stated considerations following criteria can be formulated:

Natural number  $N = 6P + 5$ ;  $P = 0, 1, 2, 3, \dots$  is a prime if and only if there is no solution for Diophantine equation

$$P = 6xy - x + y - 1; x \geq 1; y \geq 1 \quad (27)$$

Natural number  $N = 6P + 7$ ;  $P = 0, 1, 2, 3, \dots$  is a prime if and only if no one of two Diophantine equations

$$P = 6xy - x - y - 1; x \geq 1; y \geq x \quad (28)$$

$$P = 6xy + x + y - 1; x \geq 1; y \geq x \quad (29)$$

Has solution.

To prove primality criteria suppose that natural number  $N = 6P + 5$  is a prime and integer solution for equation (27) does exist. This means that there is corresponding member  $P$  in arrays  $P_1(i, j)$  or  $P_2(i, j)$ , i.e. natural number  $N = 6P + 5$  is not prime, this is a contradiction.

So natural number  $N=6P+5$ ;  $P=0, 1, 2, 3, \dots$  is a prime if and only if there is no integer solution for equation (27). The same is true for  $N=6P+7$  (28) and (29).

C++ program for primality testing of  $N$  is presented in Appendix 2.

Employment of Diophantine equation for primality testing is illustrated on the page <https://www.wolframcloud.com/objects/f8816adc-cf73-452a-90ab-29a14b763f3f>.

## 5. Twin-prime criteria

Obviously, twin primes  $N_1$  and  $N_2$  ( $N_2-N_1=2$ ) correspond to formulae

$$N_1=6P+5 \quad \text{and} \quad N_2=6P+7$$

So twin-prime criteria can be formulated as follows:

Natural numbers  $N_1=6P+5$  and  $N_2=6P+7, P=0, 1, 2, 3, \dots$  are twin primes if and only if no one of three Diophantine equation

$$P = 6xy - x + y - 1; x \geq 1; y \geq 1 \quad (28)$$

$$P = 6xy - x - y - 1; x \geq 1; y \geq x \quad (29)$$

$$P = 6xy + x + y - 1; x \geq 1; y \geq x \quad (30)$$

has solution.

## 6. Conclusions

We have shown that all composite numbers in the sequence  $S_1(P)=6P+5$ ;  $P=0, 1, 2, 3, \dots$  correspond to formula

$$S_1(P) = (5 + 6i)(7 + 6j);$$

and all composite numbers in the sequence  $S_2(P)=6P+7$ ;  $P=0, 1, 2, 3, \dots$  correspond to formulae

$$S_2(P) = (5 + 6i)(5 + 6j) \text{ or}$$

$$S_2(P) = (7 + 6i)(7 + 6j);$$

Deterministic “**matrix definition**” for prime numbers was derived:

Natural numbers that are **not** contained in arrays

$$P_1(i, j) = 6ij - i + j - 1 = \begin{pmatrix} 5 & 12 & 19 & 26 & 33 & 40 & \dots \\ 0 & 23 & 36 & 49 & 62 & 75 & \dots \\ 0 & 0 & 53 & 72 & 91 & 110 & \dots \\ 0 & 0 & 0 & 95 & 120 & 145 & \dots \\ 0 & 0 & 0 & 0 & 149 & 180 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 215 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i$$

$$P_2(i, j) = 6ij + i - j - 1 = \begin{pmatrix} 0 & 10 & 15 & 20 & 25 & 30 & \dots \\ 0 & 0 & 34 & 45 & 56 & 67 & \dots \\ 0 & 0 & 0 & 70 & 87 & 104 & \dots \\ 0 & 0 & 0 & 0 & 118 & 141 & \dots \\ 0 & 0 & 0 & 0 & 0 & 178 & \dots \\ \dots & \dots & \dots & \dots & \dots & 0 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i + 1$$

are indexes  $P$  of **all primes** in the sequence  $S_1(P) = 6P + 5$ .

Natural numbers that are **not** contained in arrays

$$P_3(i, j) = 6ij - i - j - 1 = \begin{pmatrix} 3 & 8 & 13 & 18 & 23 & 28 & \dots \\ 0 & 19 & 30 & 41 & 52 & 63 & \dots \\ 0 & 0 & 47 & 64 & 81 & 98 & \dots \\ 0 & 0 & 0 & 87 & 110 & 133 & \dots \\ 0 & 0 & 0 & 0 & 139 & 168 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 203 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i$$

$$P_4(i, j) = 6ij + i + j - 1 = \begin{pmatrix} 7 & 14 & 21 & 28 & 35 & 42 & \dots \\ 0 & 27 & 40 & 53 & 66 & 79 & \dots \\ 0 & 0 & 59 & 78 & 97 & 116 & \dots \\ 0 & 0 & 0 & 103 & 128 & 153 & \dots \\ 0 & 0 & 0 & 0 & 159 & 190 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots & 227 & \dots \end{pmatrix}; i = 1, 2, 3, 4 \dots; j \geq i$$

are indexes  $P$  of **all primes** in the sequence  $S_2(P) = 6P + 7$ .

General criteria of primality and twin-prime criteria were formulated.

C++ program for finding primes in given range  $(N1; N2)$ ,  $(N2 < 2 \cdot 10^{18}, N2 - N1 < 1\,000\,000)$  was developed and successfully tested, confirming theoretical background. For  $N = 2 \cdot 10^{18}$  run time on ordinary notebook equals 10 s.

C++ program for testing primality of given natural number  $N$  was developed and successfully tested. For  $N = 2 \cdot 10^{18}$  run time on ordinary notebook equals 7 s,

*Attachment 1*

```
#include <cstdlib>
```



```

#include <iostream>
#include <math.h>
#include <ctime>

using namespace std;
main( )
{
    /* 3S MATRIX SIEVE*/
    /*FINDING PRIMES IN THE RANGE (N1;N2)*/

    /* N1>23; N2<2 000 000 000 000 000 000; N2-N1<500 000*/

    unsigned long longint N1=111222333444555000; unsigned long longint N2
    =111222333444555600;
    unsigned long longint pr1=floor(N1/6); unsigned long longint pr2=ceil( N2/6);
    int r=84000; int R2[r]; intrm=pr2-pr1; unsigned long longint S2[r]; int r3, r4;
    int q=84000; int R1[q] ; intqm=rm; unsigned long longint S1[q] ; int q2, q1;
    for (q=1;q<qm;q++)
    R1[q] =1;
    for (r=1;r<rm;r++)
        R2[r] =1;
    unsigned long longinti, j, P1, P2, P3, P4, B, K;

    unsigned long longint i2= sqrt( pr2/6)+2;
    longlongint j1, j2;
    int l1=0;int l2=0;
    for ( i=1;i<i2;i++)
    { j2=(pr2+i+1)/( 6*i+1)+1;j1=(pr1+i+1)/( 6*i+1);
    B=5+5*( i-1); K=7+6*( i-1);
    if ( i>j1) j1=i;
    for(j=j1; j<j2;j++)
    { P1=B+K*( j-1);
    if(( P1>pr1)&&( P1<pr2))
    { q1=P1-pr1; R1[ q1] =0; } }
    j2=(pr2-i+1)/( 6*i-1)+1;j1=(pr1-i+1)/( 6*i-1);
    if (j1<1) j1=1;
    B=5+7*( i-1); K=5+6*( i-1);
    if ( i>j1-1) j1=i+1;
    for(j=j1; j<j2;j++)

```

```

{P2=B+K*(j-1);
  if(( P2>pr1)&&( P2<pr2))
  {   q2=P2-pr1; R1[ q2] =0;
  } }
  j2=(pr2+i+1)/( 6*i-1)+1;j1=(pr1+i+1)/( 6*i-1);
  B=3+5*( i-1); K=5+6*( i-1);
if ( i>j1) j1=i;
for(j=j1; P3=B+K*( j1);
if(( P3>pr1)&&( P3<pr2))
  {   r3=P3-pr1; R2[ r3]=0;
  } }
  j2=(pr2-i+1)/( 6*i+1)+1;j1=(pr1-i+1)/( 6*i+1);
  B=7+7*(i-1); K=7+6*(i-1);
if ( i>j1) j1=i;
for(j=j1; j<j2;j++)
{ P4=B+K*( j-1);
if(( P4>pr1)&&( P4<pr2))
{ r4=P4-pr1; R2[ r4] =0; } } }
cout<<"ni2="<<i2<<";pr2="<<pr2<<"  \n";
cout<<"nP=pr1+q; pr1="<<pr1<<"  \n";
  for ( q=1;q<qm;q++) { S1[q] =R1[q]*((pr1+q)*6+5); if (S1[q]%5==0) continue;
l1=l1+1;
cout<<"q="<<q<<"; Prime in S1[P]=6*P+5="<<S1[ q]<<"  \n";}
cout<<"nP=pr1+r; pr1="<<pr1<<"  \n";
for ( r=1;r<rm;r++) { S2[r] =R2[r]*((pr1+r)*6+7);if (S2[r]%5==0) continue;l2=l2+1;
cout<<"r="<<r<<"; Prime in S2[P]=6*P+7="<<S2[ r]<<"  \n";}
cout<<"l1="<<l1<<"  \n";
cout<<"l2="<<l2<<"  \n";
cout<<"number of Primes in the range (N1;N2) =l1+l2\n";
cout<<" run time (ms)=";
cout<<clock();
system("PAUSE");
return EXIT_SUCCESS;

}

```

## Attachment 2

```

#include <cstdlib>
#include <iostream>

```



```
break;}
if (i==i2-1)
cout<<"N is prime"<<"\nN="<<N<<" \n";
}}
if (mod3N==0)
cout<<"N is not prime"<<"\nN="<<N<<" \n";
cout<<clock();
system("PAUSE");
return EXIT_SUCCESS;

}
```